## Announcements

- Questions re/labs $\rightarrow$ TAs
- Yes, lab02 is due THIS week
- In case of difficulty finishing assignments $\rightarrow$ talk to TA first
- Don't worry about the diary grades. They may show up late, or as a zero if the TA bundles grades.
- First homework $\rightarrow$ next week



## A simple example: add 0.1 repeatedly 100,000 times

We know the answer to this: $\quad \sum_{k=1}^{100,000} 0.1=10,000$

This is the answer my computer gave when I used about a number scheme that had about 7 decimal digits of precision:

This is the answer my computer gave when I used about a number scheme that had about 16 decimal digits

$$
\sum_{k=1}^{100,000} 0.1=9,998.556640625
$$ of precision*:

$$
\sum_{k=1}^{100,000} 0.1=10,000.0000000188480000
$$

## $10,000.0000000188480000 \neq 10,000$

## So what happened?

We do arithmetic using decimal numbers, so this is how we also tend to define our instructions to the computer.

Almost all computers use binary numbers
... 01100011 1000...


We also prefer to get our answers in the number system we are used to (decimal)


Some things are lost in translation.
0.1
0.2
0.3
0.4
0.5
0.6
0.7
0.8
0.9
1.0

0.10000000000000001000
0.20000000000000001000
0.30000000000000004000
0.40000000000000002000
0.50000000000000000000
0.59999999999999998000
0.70000000000000007000
0.80000000000000004000
0.90000000000000002000
1.00000000000000000000

## A significant digit: one that is known to be correct and reliable



These all have three significant digits:
0.716
. 000716
716
7.16e5

These have four significant digits:
7.160e5
0.7160

This one is tricky (ambiguous):
716000
Has it been rounded?
Is it precise to the nearest 1000?
Is it precise to the nearest 1 ?


## Calculations with significant digits


$m=1.1 \mathrm{~kg}, k=4350.3142 \mathrm{~N} / \mathrm{m}$
$f_{o}=\frac{1}{2 \pi} \sqrt{\frac{k}{m}}=\underset{\uparrow}{10.008848175944461 \mathrm{~Hz}}$

A general rule:
Stated results should typically be of the same order of magnitude as the uncertainty. That is, we don't use more significant digits than we're sure about.

## Important Note:

When doing calculations on the way to an answer, you should normally be using at least one extra significant digit, and rounded at the end for the final answer.

Does this seem reasonable? The least precise data is given with two significant digits, so our answer should be also be given with two significant digits: $f_{o}=10 \mathrm{~Hz}$

# Accuracy and Precision <br> (error analysis) 

Blunders or mistakes:

- Transposed numbers
- wrong units
- incorrect decimal places

Discrepancies or disagreements

- The world is flat. No, its round. (actually, its an oblate spheroid)



## Uncertainty


systematic error (bias)
random errors


Not precise


## Absolute and relative error

| F | C |
| :---: | :---: |
| $120=$ | 50 |
| 100 | 40 |
| 80 | 30 |
| 60 | 20 |
|  | 10 |
|  | 0 |
| 20 | -10 |
| 0 | -20 |
| $-20$ | -30 |
| -40 三 | -40 |

Truth: $\quad T_{t}=1^{\circ} \mathrm{C}$
Measurement: $T_{a}=0^{\circ} \mathrm{C}$

True error: $\quad \epsilon=\left|T_{t}-T_{a}\right|$
Relative error: $\quad \eta=\frac{\left|T_{t}-T_{a}\right|}{\left|T_{t}\right|}=\frac{\epsilon}{\left|T_{t}\right|}$
Percent error: $\quad \frac{\epsilon}{\left|T_{t}\right|} \times 100$
Problem: Do we ever know the truth?

Questions for further thought: What is the relative error in the example above if the truth is $10^{\circ}$ ? Or $100^{\circ}$ ?

What is the relative error if the truth is $0^{\circ}$ ?

## An example (where we don't know the truth)

Binomial expansion
$(1+x)^{a}=1+a x+\frac{a(a-1)}{2!} x^{2}+\frac{a(a-1)(a-2)}{3!} x^{3}+\frac{a(a-1)(a-2)(a-3)}{4!} x^{3}+\cdots$
$|x|<1, a \neq 0$, real

Compute $1.4^{3.1}$ to within $10 \%$ approximate error, using as few terms as possible

First
approx.: $(1+0.4)^{3.1}=1+3.1 * 0.4=2.24$
$\begin{aligned} & \text { Second } \\ & \text { approx.: }\end{aligned}(1+0.4)^{3.1}=1+3.1 * 0.4+\frac{3.1(3.1-1)}{2!} 0.4^{2}=2.76$
$|2.24-2.76| / 2.76=0.19$

Third approx.:

$$
(1+0.4)^{3.1}=1+3.1 * 0.4+\frac{3.1(3.1-1)}{2!} 0.4^{2}+\frac{3.1(3.1-1)(3.1-2)}{3!} 0.4^{3}=2.84
$$

$$
|2.76-2.84| / 2.84=0.028<10 \%
$$

## Numbers in Computers



Base-10:

$$
247 \rightarrow 2^{*} 10^{2}+4^{*} 10+7^{*} 10^{0}
$$



## Numbers in Computers

signed
8-bit (1-byte) binary numbers can represent these integers:

16-bit (2-byte) binary
numbers can represent these integers:

24-bit (3-byte) binary numbers can represent these integers:
32-bit (4-byte) binary numbers can represent these integers:

64-bit (8-byte) binary numbers can represent these integers:
$-128 \rightarrow 127$
C: char; Fortran:
INTEGER*1
$-32,768 \rightarrow 32,767$
C: short int
Fortran: INTEGER*2

Not used
$-8,388,608 \rightarrow 8,388,607$
$0 \rightarrow 16,777,215$

Fortran: INTEGER, INTEGER*4

$$
0 \rightarrow 65,537
$$

unsigned
$0 \rightarrow 255$

$$
\begin{array}{cc}
\sim+-10^{9} & 0 \rightarrow 4,294,967,295 \\
\mathrm{C}: \text { int } & \text { C: unsigned int }
\end{array}
$$

$$
\begin{array}{cc}
\sim+-10^{18} & 0 \rightarrow \\
\text { C: long int } & 1.84467440737096 \mathrm{e}+19 \\
\text { Fortran: INTEGER*8 } & \text { C: unsigned long int }
\end{array}
$$

## Efficiency of Binary Numbers



To access these photos, I need 8 unique numbers using base-10.

But, what if we address these photos using a binary representation of their address?


That's not bad for 8 locations. What is the efficiency for 64,000 locations?

## Numbers in Computers

Base-10:

$$
247.13 \rightarrow 2^{*} 10^{2}+4 * 10+7^{*} 10^{0}+1^{*} 10^{-1}+3^{*} 10^{-2}
$$



## Octal/Hexadecimal/Base 256

Base-8 (octal):
Group 3 bits in binary: $110010101010 \rightarrow 06252$ or 06252
No digit larger than 7!

Base-16 (hexadecimal):
Group 4 bits in binary: $110010100010 \rightarrow 0 x D A 2$
The digits are now: 0123456789 ABCDEF

Base-256 (IP addresses):

```
Configure IPv4: Using DHCP \hat{ }
    IPv4 Address: 10.21.96.14
    Subnet Mask: 255.255.0.0 DHCP Cl
    Router: 10.21.0.1
```

$\rightarrow$ Integer numbers on a computer have a limited range!
$\rightarrow$ Integer numbers have a constant true error (0.5)!

What can we do if we need larger numbers?

## Floating Point Numbers

A floating point number:


Questions:
Can I represent 3.928 with the 'boxes' (or bits) shown above? What is the precision of this floating point number?

## Anatomy of a 32-bit (a.k.a. single) Floating Point Number

##  <br> sign <br> bits <br>  <br> 7 bits for representing the exponent <br> 23 bits for representing the mantissa

$10^{-39} \rightarrow 10^{38}, \sim 7$ digits of precision

64-bit ('double') floating point number:
11-bit exponent
52-bit mantissa
2 sign

$$
10^{-308} \rightarrow 10^{308}, \sim 16 \text { digits of precision }
$$

## 64 bit (8-byte) floating point number

IEEE 754 standard, used internally by MATLAB


The real value assumed by a given 64-bit double-precision datum with a given biased exponent $e$ and a 52 -bit fraction is

$$
(-1)^{\mathrm{sign}}\left(1 . b_{51} b_{50} \ldots b_{0}\right)_{2} \times 2^{e-1023}
$$

or

$$
(-1)^{\text {sign }}\left(1+\sum_{i=1}^{52} b_{52-i} 2^{-i}\right) \times 2^{e-1023}
$$

Between $2^{52}=4,503,599,627,370,496$ and $2^{53}=9,007,199,254,740,992$ the representable numbers are exactly the integers. For the next range, from $2^{53}$ to $2^{54}$, everything is multiplied by 2 , so the representable numbers are the even ones, etc. Conversely, for the previous range from $2^{51}$ to $2^{52}$, the spacing is 0.5 , etc.

## MATLAB example: very small numbers (de-normalized numbers, gradual underflow


$\rightarrow$ Floating point numbers have a much larger range than integers with the same storage requirement
$\rightarrow$ Floating point numbers have a (more or less) constant relative error (precision)
$\rightarrow$ Only a very limited subset of real numbers can be represented on a computer

# Sick cases (usually coding error) 

## 1/0, 2e222^2 $\rightarrow$ Floating point overflow $\rightarrow$ Inf

## $2 \mathrm{e}-222^{\wedge} 2 \rightarrow$ Floating point underflow $\rightarrow 0$

$0 / 0 \rightarrow$ Makes no sense $\rightarrow \mathrm{NaN}$

## MATLAB example

```
clear all; close all;
% this is a simple matlab script to examine the percent error
% in very small numbers using MATLAB's native doubl precision
% number scheme.
8
% The script examines the difference in the ratio
%
% 1.1 x 10^-exponent
8
% }\quad1.0\times1\mp@subsup{0}{}{\wedge}\mathrm{ -exponent
%
exponent = 300:330;
for i = 1:length (exponent)
    num = 1.1*10^(-exponent(i));
    den = 1.0*10^(-exponent(i));
    x(i) = num/den;
end
perError = 100*abs(x-1.1)/1.1; % this is the percent error
% plot the result in a linear plot (is this hard to see?)
subplot (211)
plot(exponent,perError,'o')
xlabel('exponent')
ylabel('% Error')
% plot the result in a log plot
subplot(212)
semilogy (exponent, perError,'0')
xlabel ('exponent')
ylabel('% Error')
```

Results


## Question:

The MATLAB code uses exponents as high as 330 ( $10^{\wedge}-330$ ). Why don't we see this in our plot? What does MATLAB give as a \% error for these high exponents?

## Quantization Error

Limitations in precision leads to truncation or rounding Example: audio sampling (16 bit ADC typically)







## Return to the first example:

This is the answer my computer gave when I used about a number scheme that had about 7 decimal digits of precision:

$$
\sum_{k=1}^{100,000} 0.1=9,998.556640625
$$

The individual error in precision is small, but because were doing a large number of computations that depend on the results of earlier ones, the error grows large:


## How to represent text?

Each character is one byte $\rightarrow$ ascii (American Standard Code for Information
Interchange) table. 0-31 are control characters, 128-255 are extras, some are not
printable.

| Dec |  | x Oct | Char |  | De |  | Oct | Html Chr | Dec |  | x Oct |  | chr |  | HX 0 | Html Ch |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 |  | 000 | NUL | (null) | 32 | 20 | 040 | \&\#32: space | 64 | 440 | 100 | \&\#64; | 0 | 96 | 60140 | \&\#96; |  |
| 1 | 1 | 001 | 50H | (start of heading) | 33 | 21 | 041 | \&\#33; | 65 | 51 | 101 | \&\#65; | A | 97 | 61141 | \&\#97; | a |
| 2 | 2 | 002 | STX | (start of text) | 34 | 22 | 042 | \&\#34; | 66 | 642 | 102 | \&\#66; | B | 98 | 62142 | \&\#98; | b |
| 3 | 3 | 003 | ETX | (end of text) | 35 | 23 | 043 | \&\#35; \# | 67 | 743 | 103 | \&\#67; | C | 99 | 63143 | \&\#99; | c |
| 4 | 4 | 004 | E0T | (end of transmission) | 36 | 24 | 044 | \&\#36; | 68 | 44 | 104 | \& 68 $^{\text {\% }}$ | D | 100 | 64144 | \&\#100; | d |
| 5 | 5 | 005 | ENQ | (enquiry) | 37 | 25 | 045 | \&\#37; \% | 69 | 45 | 105 | \&\#69; | E | 101 | 65145 | \&\#101; | e |
| 6 | 6 | 006 | ACK | (acknowledge) | 38 | 26 | 046 | \&\#38; | 70 | 46 | 106 | \&\#70; | F | 102 | 66146 | \&\#102; | f |
| 7 | 7 | 007 | BEL | (bell) | 39 | 27 | 047 | \&\#39; | 71 | 147 | 107 | \&\#71; | G | 103 | $67 \quad 147$ | \&\#103; | 9 |
| 8 | 8 | 010 | BS | (backspace) | 40 | 28 | 050 | \&\#40; | 72 | 48 | 110 | \&\#72; | H | 104 | 68150 | \&\#104; | h |
| 9 | 9 | 011 | TAB | (horizontal tab) | 41 | 29 | 051 | \&\#41; | 73 | 49 | 111 | \&\#73; | I | 105 | 69151 | \&\#105; | i |
| 10 | A | 012 | LF | (NL line feed, new line) | 42 | 2A | 052 | \&\#42; | 74 | 4 4A | 112 | \&\#74; | J | 106 | 6A 152 | \&\#106; | j |
| 11 | B | 013 | VT | (vertical tab) | 43 | 2B | 053 | \&\#43 | 75 | 54 B | 113 | \&\#75; | K | 107 | 6B 153 | \&\#107 | k |
| 12 | C | 014 | FF | (NP form feed, new page) | 44 | 2C | 054 | \&\#44; | 76 | 6 4C | 114 | \&\#76; | L | 108 | 6C 154 | \&\#108; | 1 |
| 13 | D | 015 | CR | (carriage return) | 45 | 2D | 055 | \&\#45; | 77 | 7 4D | 115 | \&\#77: | M | 109 | 6D 155 | \&\#109; | mil |
| 14 | E | 016 | S0 | (shift out) | 46 | 2E | 056 | \&\#46; | 78 | $8 \mathrm{4E}$ | 116 | \&\#78; | N | 110 | 6 E 156 | \&\#110; | n |
| 15 | F | 017 | SI | (shift in) | 47 | 2 F | 057 | \&\#47; | 79 | 4 F | 117 | \&\#79; | 0 | 111 | 6 F 157 | \&\#111 | 0 |
| 16 | 10 | 020 | DLE | (data link escape) | 48 | 30 | 060 | \&\#48; 0 | 80 | 50 | 120 | \&\#80; | P | 112 | 70160 | \&\#112; | p |
| 17 | 11 | 021 | DC1 | (device control 1) | 49 | 31 | 061 | \&\#49; | 81 | 151 | 121 | \&\#81; | Q | 113 | 71161 | \&\#113; | q |
| 18 | 12 | 022 | DC2 | (device control 2) | 50 | 32 | 062 | \&\#50; 2 | 82 | 52 | 122 | \&\#82; | R | 114 | 72162 | \&\#114; | r |
| 19 | 13 | 023 | DC3 | (device control 3) | 51 | 33 | 063 | \&\#51:3 | 83 | 53 | 123 | \&\#83; | 5 | 115 | 73163 | \&\#115; | 3 |
| 20 | 14 | 024 | DC4 | (device control 4) | 52 | 34 | 064 | \&\#52; 4 | 84 | 454 | 124 | \&\#84; | T | 116 | 74164 | \&\#116; | t |
| 21 | 15 | 025 | NAK | (negative acknowledge) | 53 | 35 | 065 | \&\#53; 5 | 85 | 55 | 125 | \&\#85; | U | 117 | 75165 | \&\#117: | u |
| 22 | 16 | 026 | SYN | (synchronous idle) | 54 | 36 | 066 | \&\#54; 6 | 86 | 56 | 126 | \&\#86; | V | 118 | 76166 | \&\#118; | v |
| 23 | 17 | 027 | ETB | (end of trans. block) | 55 | 37 | 067 | \&\#55; 7 | 87 | 75 | 127 | \&\#87: | W | 119 | 77167 | \&\#119; | W |
| 24 | 18 | 030 | CAN | (cancel) | 56 | 38 | 070 | \&\#56; 8 | 88 | 88 | 130 | \&\#88; | X | 120 | 78170 | \&\#120; | X |
| 25 | 19 | 031 | EM | (end of medium) | 57 | 39 | 071 | \&\#57; 9 | 89 | 59 | 131 | \&\#89; | Y | 121 | 79171 | \&\#121; | Y |
| 26 | 1 A | 032 | SUB | (substitute) | 58 | 3A | 072 | \&\#58; | 90 | 5A | 132 | \&\#90: | Z | 122 | 7A 172 | \&\#122; | z |
| 27 | 1B | 033 | ESC | (escape) | 59 | 3B | 073 | \&\#59; | 91 | 1 5B | 133 | \&\#91; | [ | 123 | 7B 173 | \&\#123; |  |
| 28 | 1 C | 034 | FS | (file separator) | 60 | 3C | 074 | \&\#60; < | 92 | 2 5C | 134 | \&\#92; | \} | 124 | 7C 174 | \&\#124; | 1 |
| 29 | 1D | 035 | GS | (group separator) | 61 | 3D | 075 | \&\#61 | 93 | 3 5D | 135 | \&\#93; | 1 | 125 | 7D 175 | \&\#125; |  |
| 30 | 1 E | 036 | RS | (record separator) | 62 | 3E | 076 | \&\#62; > | 94 | 4 5E | 136 | \&\#94; | 人 |  | 7E 176 | \&\#126; |  |
| 31 | 1 F | 037 | US | (unit separator) | 63 | 3 F | 077 | \&\#63; ? | 95 | 5 FF | 137 | \&\#95; | - | 127 | 7F 177 | \&\#127; | DEL |

Source: www.Lookup Tables.com

## How to represent text and other stuff?

One line of text $\rightarrow$ 'text'<lf> or 'text<cr><|f>
End of file $\rightarrow$ <ctrl>Z (26) only Microsoft
How to store a picture $\rightarrow 3$ bytes per pixel rgbrgbrgb..... Then compress (jpg)
How to store sound $\rightarrow$ 16bit (2byte) samples at $44,200 \mathrm{~Hz} \rightarrow 5304000$ bytes/minute $\rightarrow \sim 100$ minutes on a CDROM (wav file).

## Number Representations in Computers

character
'char'
A 1-byte individual character
ASCII Characters:
!"\#\$\% \& ( ) * + , - / 01
$23456789: ;<=>$ ? @
BCDEFGHIJKLMNOP
QRSTUVWXYZ[\]^_` abcdefghljklmnopq rstuvwxyz\{|\}~

In the MATLAB command window:
>> char(32:126)
integer
int

An integer value
Either signed:
...-5,-4,-3,-2,-1,0,1,2,3,4,...
or unsigned:

$$
0,1,2,3,4,5, \ldots
$$

The number of values that can be represented depends on the number of bytes:

1 byte (unsigned):
$0 \rightarrow 255$
2 byte (short, unsigned): $0 \rightarrow 65,535$
4 byte (long, unsigned):
$0 \rightarrow 4,294,967,296$
floating point point

A floating point value
Single (4 bytes)

- 6-9 significant decimal points
- max value is $\sim 3 \times 10^{38}$

Double (8 bytes)

- 15-17 significant decimal
points
- max value is $\sim 1 \times 10^{308}$

Quadruple (16 bytes)

Reference: Kernighan, B. W., and D. M. Ritchie, The C Programming Language, Second Edition, Prentice-Hall, Inc., 1988.

## Unit Systems

Table 1. SI base units
http://physics.nist.gov/Pubs/SP330/sp330.pdf

| Base quantity |  | SI base unit |  |
| :---: | :---: | :---: | :---: |
| Name | Symbol | Name | Symbol |
| length | $l, x, r$, etc. | meter | m |
| mass | $m$ | kilogram | kg |
| time, duration | $t$ | second | A |
| electric current | $I, i$ | ampere | A |
| thermodynamic temperature | $T$ | kelvin | K |
| amount of substance | $n$ | mole | mol |
| luminous intensity | $I_{\mathrm{v}}$ | candela | cd |

Air Canada Flight 143: "Gimli Glider"


## TABLE 1.1 SI Units

| Quantity | Name of unit | Symbol | Equivalent |
| :--- | :--- | :--- | :--- |
| Length | Meter | m |  |
| Mass | Kilogram | kg |  |
| Time | Second | s |  |
| Temperature | Kelvin | K |  |
| Frequency | Hertz | Hz | $\mathrm{s}^{-1}$ |
| Force | Newton | N | $\mathrm{kg} \mathrm{ms}^{-2}$ |
| Pressure | Pascal | Pa | $\mathrm{N} \mathrm{m}^{-2}$ |
| Energy | Joule | J | $\mathrm{N} \mathrm{m}^{\text {Power }}$ |

TABLE 1.2 Common Prefixes

| Prefix | Symbol | Multiple |
| :--- | :--- | :--- |
| Mega | M | $10^{6}$ |
| Kilo | k | $10^{3}$ |
| Deci | d | $10^{-1}$ |
| Centi | c | $10^{-2}$ |
| Milli | m | $10^{-3}$ |
| Micro | $\mu$ | $10^{-6}$ |

## Unit Systems

European scientists estimate that 100 kg of is wasted per person per year.


This is probably accurate to the nearest 50 kg or so, based on the available information in this sentence.

European scientists estimate that 220 lbs of food is wasted per person per year.


This is probably accurate to the nearest 5 lbs or so, based on the available information in this sentence.


Take home message: pay attention to what you are asking your computer, and to what your computer is give you for an answer.

- Significant digits
- Accuracy and precision
- Number systems
- Quantization error

